

Scalable Verification of Systems with Cryptography

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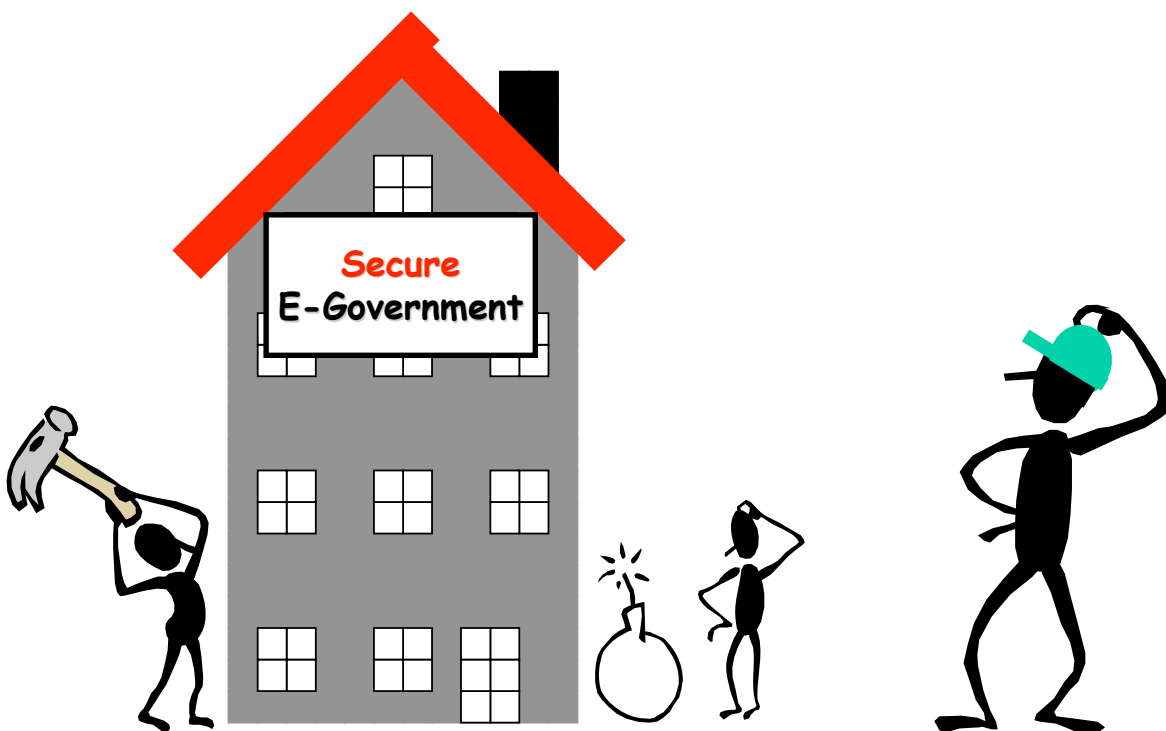
Joint work mainly with Michael Backes (Univ. Saarbrücken)
and Michael Waidner (IBM Research & SWG)



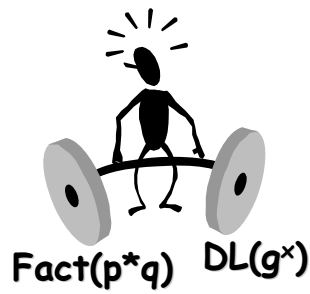
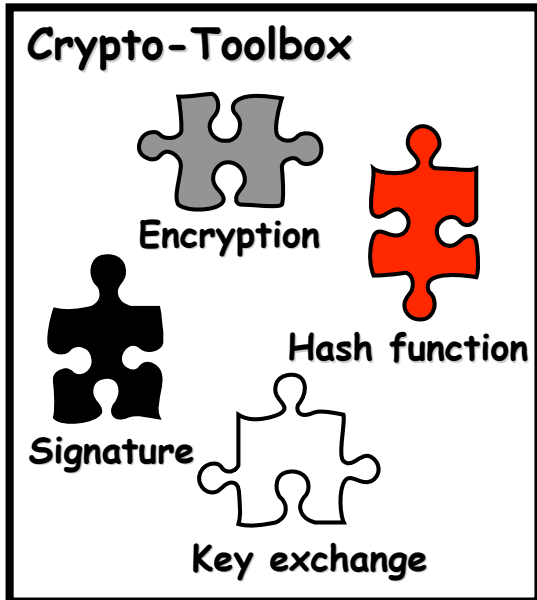
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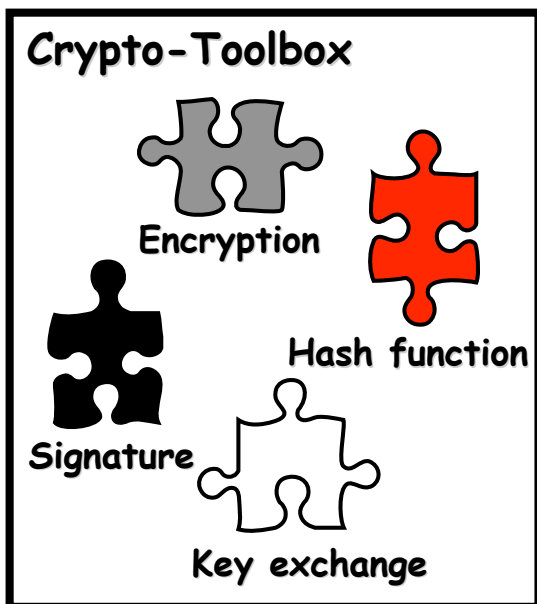
Building **Secure** Systems



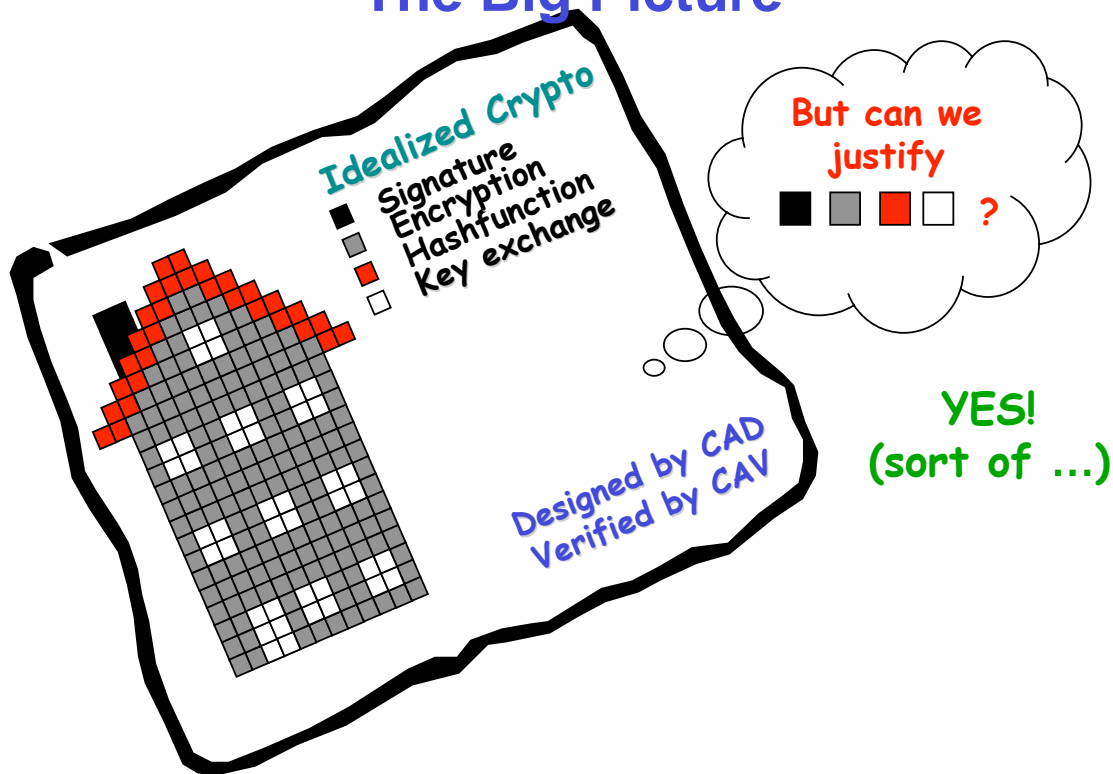
Cryptography: The Details



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Prior Automated Crypto Protocol Proofs: The Big Picture



E.g.: Secure Channels like SSL (with mutual authentication)

- If you use them in a larger system, what would you assume about them, or how would you model them?
- E.g., as “ideal secure channel”



- E.g., as a primitive in π -calculus etc.

Secure Channels, ctd.

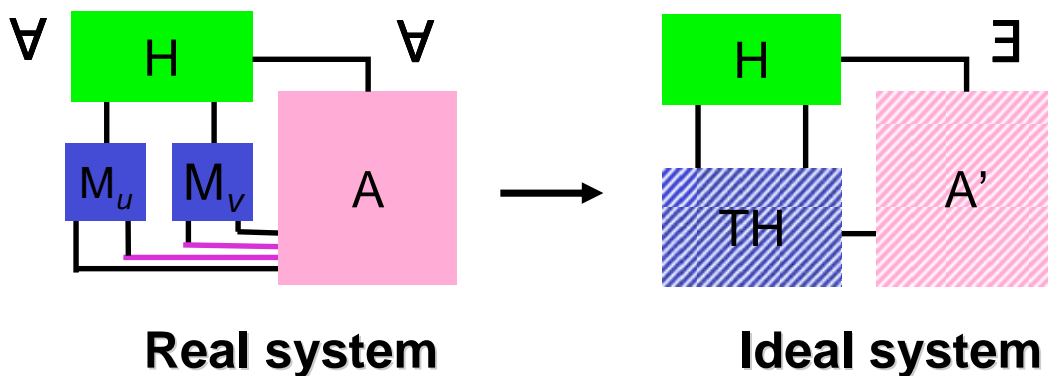


- How correct is this compared with actual SSL?
- Not bad, but not quite correct:

- Computational assumptions and error probabilities from crypto } Always very similar ⇒ make part of semantics (“fulfillment” relation)
- Message length and traffic pattern leak } Special ⇒ extend specification
- No availability } Rather general ⇒ can just be in asynchronous model

Reactive Simulatability (RSIM)

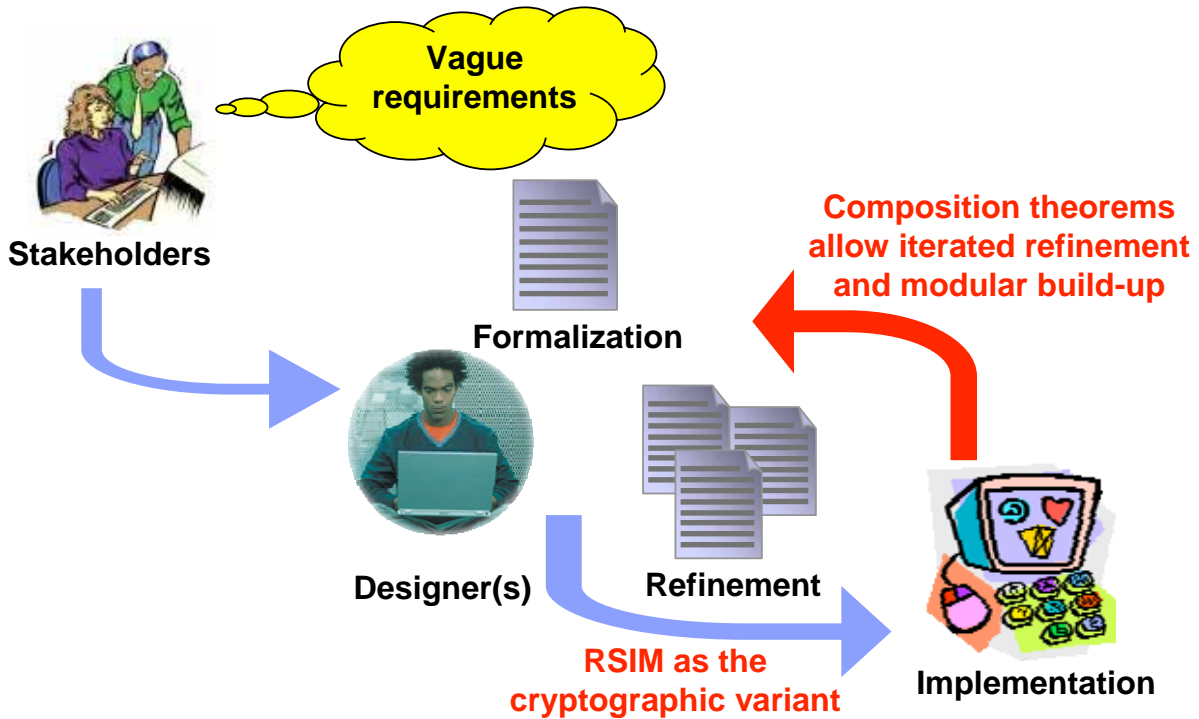
Here “General RSIM” variant



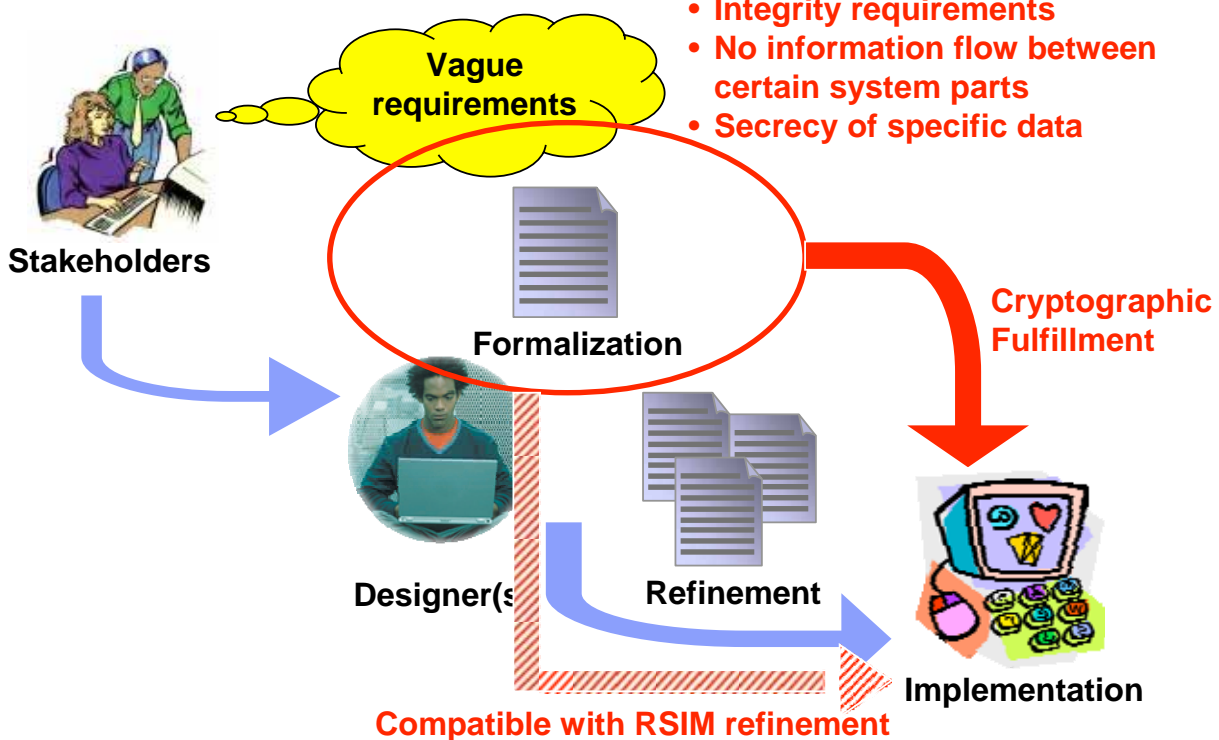
$$\text{view}_{\text{real}}(\mathbf{H}) \approx \text{view}_{\text{ideal}}(\mathbf{H})$$

Indistinguishability of random variables

RSIM in Overall Design Process



Treating Properties Cryptographically



Recent Work

- **Extended prior results for “Dolev-Yao models” – specific term-algebra abstractions widely used in verification community**
- **Impossibility results for certain Dolev-Yao model variants**
- **BPW-Dolev-Yao model in Isabelle/HOL (with Ch. Sprenger and D. Basin)**
- **Attempt to apply to real-world Web Services**